

FS \rightarrow BFS

$$A\vec{x} = \vec{b}$$

$$\vec{x} \geq 0$$

FS and

invertible

$$A\vec{x} = \begin{bmatrix} B & R \end{bmatrix} \begin{bmatrix} \vec{x}_B \\ 0 \end{bmatrix} = \vec{b}$$

$$\vec{x}_B \geq 0$$

LPP $\max \vec{c}^T \vec{x}$

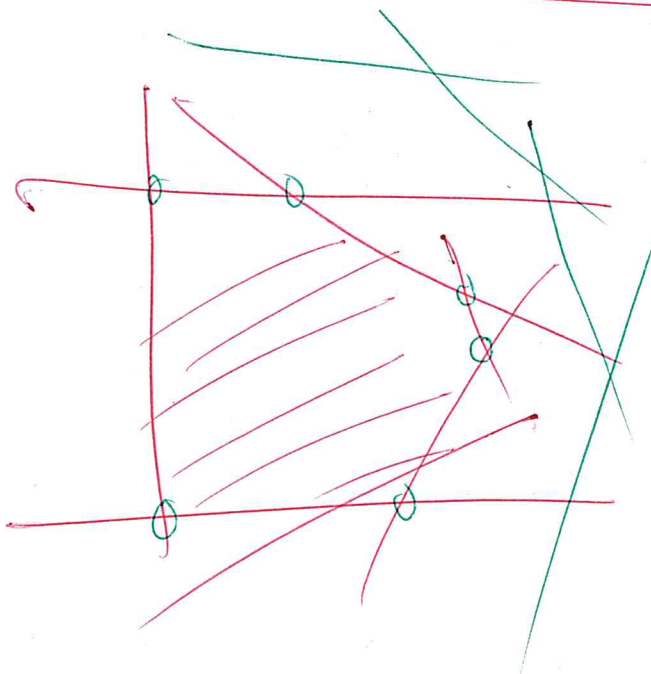
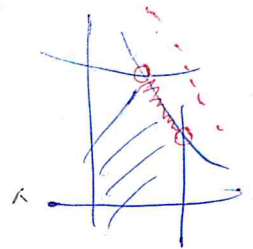
$$\begin{cases} A\vec{x} = \vec{b} \\ \vec{x} \geq 0 \end{cases}$$

rank(A) = m

A m-by-n

BFS \rightarrow corner pt of FR

optimal solution occur at BFS



Thm 2.1

BFS \rightarrow Extreme pt of FR.

Pf Let $\vec{x} \in \text{BFS}$ $\vec{x} = \begin{pmatrix} \vec{x}_B \\ \vec{0} \end{pmatrix}$

(i) $[B | R] \begin{bmatrix} \vec{x}_B \\ \vec{0} \end{bmatrix} = \vec{b}$ B invertible,

(ii) $\vec{x}_B \geq \vec{0}$

Pf by contradiction assume not an extreme pt.

$\exists \vec{x}_1 \neq \vec{x}_2 \in \text{FR}$ st.

$\vec{x} = \lambda \vec{x}_1 + (1-\lambda) \vec{x}_2, \lambda \in (0,1)$

$\vec{x}_1, \vec{x}_2 \in \text{FS}$: (1) $[B | R] \begin{pmatrix} \vec{u}_1 \\ \vec{v}_1 \end{pmatrix} = \vec{b}$ (2) $[B | R] \begin{pmatrix} \vec{u}_2 \\ \vec{v}_2 \end{pmatrix} = \vec{b}$ (1)

$\underbrace{\vec{u}_1}_{\vec{x}_1}, \vec{v}_1 \geq \vec{0}$ $\underbrace{\vec{u}_2}_{\vec{x}_2}, \vec{v}_2 \geq \vec{0}$

$\begin{pmatrix} \vec{x}_B \\ \vec{0} \end{pmatrix} = \lambda \begin{pmatrix} \vec{u}_1 \\ \vec{v}_1 \end{pmatrix} + (1-\lambda) \begin{pmatrix} \vec{u}_2 \\ \vec{v}_2 \end{pmatrix} \quad \lambda \in (0,1)$

$\begin{cases} \vec{0} = \lambda \vec{v}_1 + (1-\lambda) \vec{v}_2 & \lambda \in (0,1) \\ \Rightarrow \vec{v}_1 = \vec{v}_2 = \vec{0} \end{cases}$ (2)

(2) \Rightarrow (1) $B \vec{u}_1 = \vec{b} \quad B \vec{u}_2 = \vec{b}$

$\Rightarrow \vec{u}_1 = B^{-1} \vec{b} = \vec{u}_2$ (3)

(2) & (3) $\Rightarrow \vec{x}_1 = \vec{x}_2$ Contradiction

Thm 2.4 Extrempt of FR \Rightarrow BFS

pt: $\vec{x}_0 = \begin{pmatrix} x_1 \\ \vdots \\ x_r \\ 0 \\ \vdots \end{pmatrix}$ r can be equal to n

ex pt of FR. (FS)

$$\begin{cases} A\vec{x}_0 = \vec{b} \\ \vec{x}_0 \geq 0 \end{cases} \Leftrightarrow \begin{cases} [\vec{a}_1, \vec{a}_2, \dots, \vec{a}_r, \vec{a}_{r+1}, \dots, \vec{a}_m] \begin{pmatrix} x_1 \\ \vdots \\ x_r \\ 0 \\ \vdots \end{pmatrix} = \vec{b} \\ x_1, \dots, x_r \geq 0 \end{cases} \quad (1)$$

Case 1: ~~If~~ $\{\vec{a}_1, \dots, \vec{a}_r\}$ is li

Case 1 ~~is~~ \dots is ld

Case 1: $r \leq m = \text{rank}(A)$

~~Case~~ \exists another $m-r$ columns of A st

$\{\vec{a}_1, \dots, \vec{a}_r, \vec{a}_{r+1}, \dots, \vec{a}_m\}$ is li.

$$\left[\underbrace{\vec{a}_1 \dots \vec{a}_r}_{\text{invertible}} \mid \vec{a}_{r+1} \dots \vec{a}_m \right] \begin{pmatrix} x_1 \\ \vdots \\ x_r \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \vec{b}$$

$$[B \mid R] \begin{pmatrix} x_1 \\ \vdots \\ x_r \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \vec{b} \quad (\text{degenerate}) \Rightarrow \text{BFS}$$

Case 2 $\{\vec{a}_1, \dots, \vec{a}_r\}$ l.d. (Can never happen)

Proof by contradiction:

Suppose $\{\vec{a}_1, \dots, \vec{a}_r\}$ l.d.

$\exists \alpha_1, \dots, \alpha_r$ not all zero st

$$\alpha_1 \vec{a}_1 + \dots + \alpha_r \vec{a}_r = \vec{0} \Leftrightarrow$$

$$[\vec{a}_1, \vec{a}_2, \dots] \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_r \end{bmatrix} = \vec{0}$$

$$\epsilon \equiv \min_{\substack{1 \leq i \leq r \\ |\alpha_i| \neq 0}} \left\{ \frac{x_i}{|\alpha_i|} \right\} > 0$$

$$\Leftrightarrow A \vec{\alpha} = \vec{0}$$

$$\vec{x}_{\pm} = \vec{x}_0 \pm \epsilon \vec{\alpha}$$

$$\begin{aligned} \vec{x}_+ &= \vec{x}_0 + \epsilon \vec{\alpha} \\ \vec{x}_- &= \vec{x}_0 - \epsilon \vec{\alpha} \end{aligned}$$

$$(i) \quad \vec{x}_0 = \frac{1}{2} \vec{x}_+ + \frac{1}{2} \vec{x}_-$$

\uparrow ?? \uparrow ??
 FR FR

Cannot be an expt of FR

$$\frac{x_i}{|\alpha_i|} \geq \epsilon, \forall i \text{ } |\alpha_i| \neq 0$$

$$\Rightarrow x_i \geq \epsilon |\alpha_i| \quad \forall i$$

(ii) \vec{x}_+, \vec{x}_- feasible: $\vec{x}_+ \geq \vec{0}$
 $\vec{x}_- \geq \vec{0}$

$$\Rightarrow \underline{x_i \pm \epsilon \alpha_i \geq 0 \quad \forall i}$$

(iii) \vec{x}_+, \vec{x}_- solution $A \vec{x}_+ = \vec{b} \quad ??$
 $A \vec{x}_- = \vec{b} \quad ??$

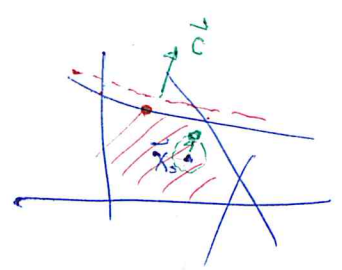
$$\begin{aligned} A \vec{x}_+ &= A(\vec{x}_0 + \epsilon \vec{\alpha}) \\ &= A \vec{x}_0 + \epsilon A \vec{\alpha} \\ &= \vec{b} + \vec{0} = \vec{b} \end{aligned}$$

Thm 2.5 Optimal Solution of LPP

occurs at a pt of FR

- pt: $\vec{x}_0 \in FR$ (i) $A\vec{x}_0 = \vec{b}$
 (ii) $\vec{x}_0 \geq \vec{0}$
 (iii) $\vec{c}^T \vec{x}_0 \geq \vec{c}^T \vec{x} \quad \forall \vec{x} \in FR$

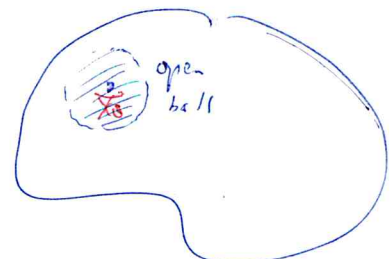
Step 1 \vec{x}_0 cannot be an interior pt



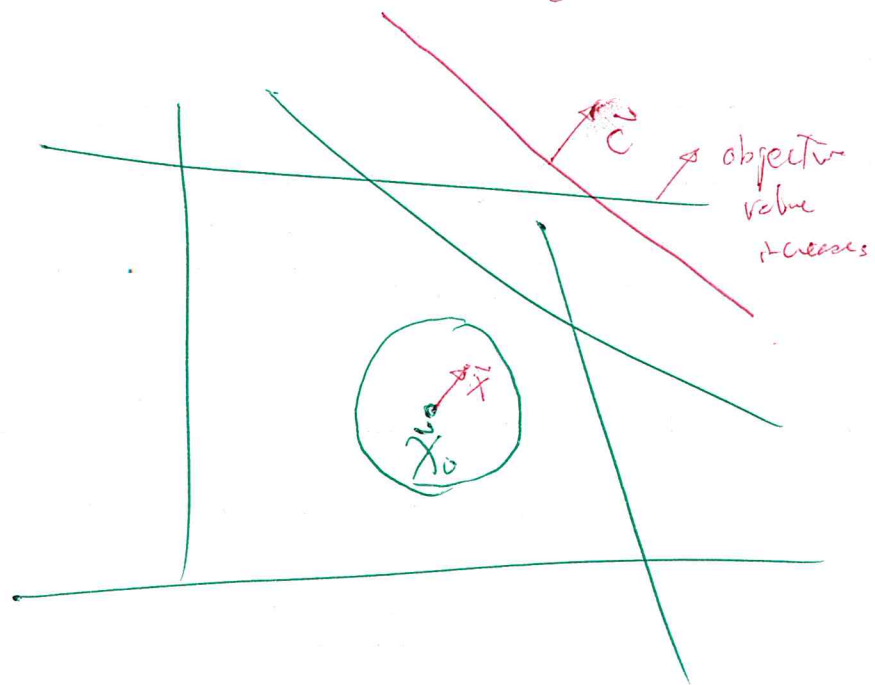
Step 2 \vec{x}_0 boundary pt $\rightarrow \exists$ corner pt that is optimal

Step 1: \vec{x}_0 is interior pt

pf by contradiction



$\exists B_\epsilon(\vec{x}_0) \subseteq FR \quad \epsilon > 0$
 $B_\epsilon(\vec{x}_0) = \{ \vec{x} \mid |\vec{x} - \vec{x}_0| < \epsilon \}$



$$\vec{x} = \vec{x}_0 + \frac{\epsilon}{2} \frac{\vec{c}}{|\vec{c}|} \in \text{FR}$$

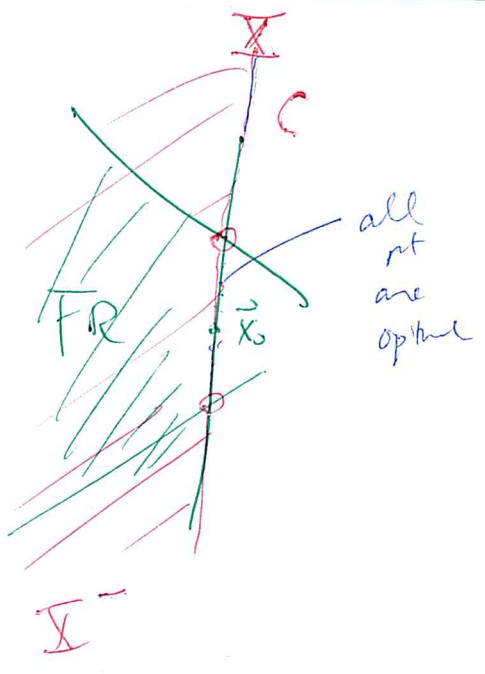
$$\vec{c}^T \vec{x} = \underbrace{\vec{c}^T \vec{x}_0}_{\phi = \text{original optimal value}} + \underbrace{\frac{\epsilon}{2} \frac{\vec{c}^T \vec{c}}{|\vec{c}|}}_{> 0}$$

a candidate

Step 2 \vec{x}_0 is boundary pt

Define $\Sigma = \{ \vec{x} \mid \vec{c}^T \vec{x} = \underbrace{\vec{c}^T \vec{x}_0}_{\text{optimal value}} \}$

- (i) $\vec{x}_0 \in \Sigma$ support of \vec{x}_0
- (ii) $\Sigma^- = \{ \vec{x} \mid \vec{c}^T \vec{x} \leq \vec{c}^T \vec{x}_0 \}$



$$\vec{x} \in \text{FR} \subseteq \Sigma^-$$

$\Rightarrow \Sigma$ is a support hyperplane of FR at \vec{x}_0

Thm 1.5 $\Rightarrow \exists$ ex pt of FR on Σ .

\Rightarrow that ex pt has to be optimal
(\because every pt on Σ is optimal)

#

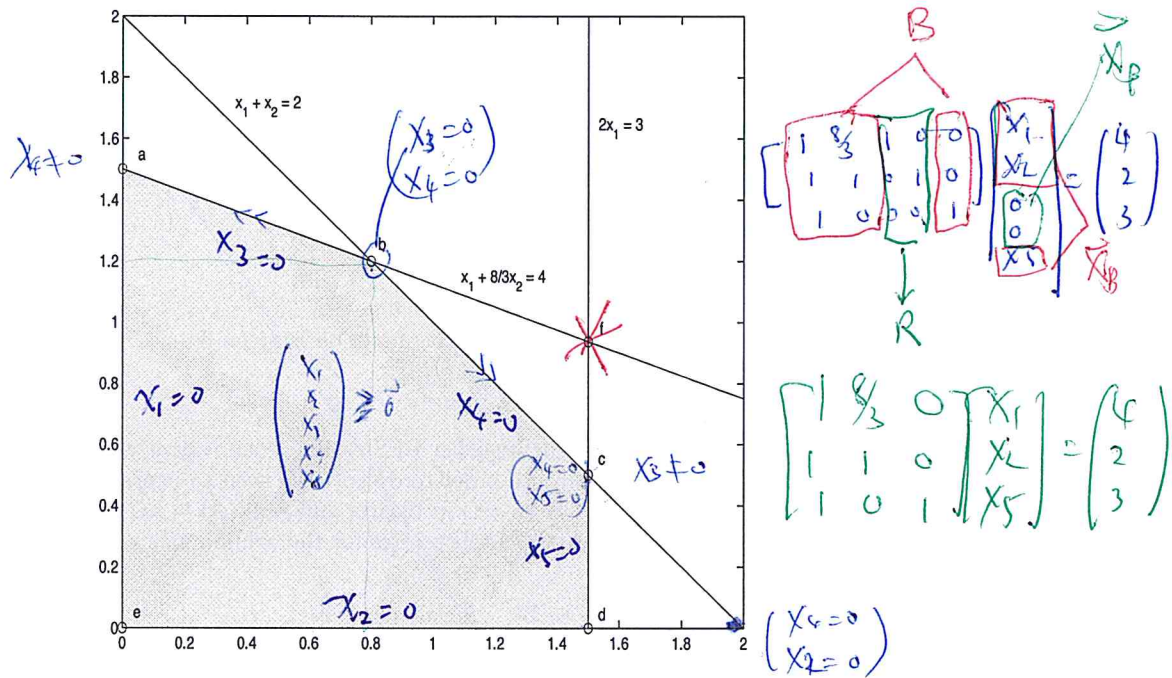


Figure 2.1. There are 5 extreme points by inspection $\{a, b, c, d, e\}$.

We get the point $[0, \frac{3}{2}, 0, \frac{1}{2}, 3]$. From Figure 2.3, we see that this point corresponds to the extreme point a of the convex polyhedron K defined by (2.19), (2.20), (2.21), (2.22). The equations $x_1 = 0$ and $x_3 = 0$ are called the *binding equations* of the extreme point a .

We note that not all basic solutions are feasible. In fact there is a maximum total of $\binom{5}{3} = \binom{5}{2} = 10$ basic solutions and here we have only five extreme points. The extreme points of the region K and their corresponding binding equations are given in the following table.

extreme point	a	b	c	d	e
set to zero	x_1 x_3	x_3 x_4	x_4 x_5	x_2 x_5	x_1 x_2

If we set x_3 and x_5 to zero, the basic solution we get is $[\frac{3}{2}, \frac{15}{16}, 0, \frac{7}{16}, 0]$. Hence it is not a basic *feasible* solution and therefore does not correspond to any one of the extreme point in K . In fact, it is given by the point f in the above figure.

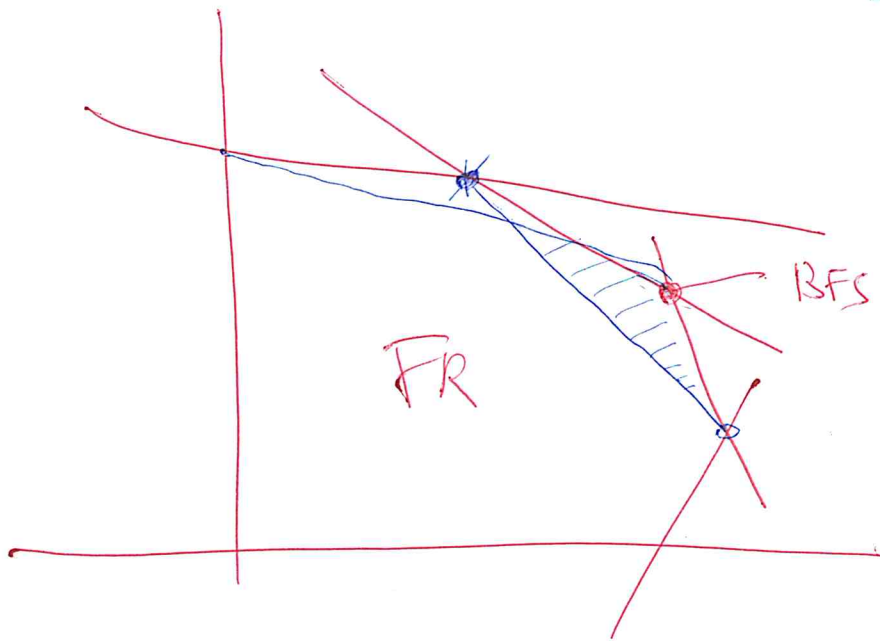
2.3 Improving a Basic Feasible Solution

Let the LPP be

$$\begin{aligned} \max \quad & z = c^T x \\ \text{subject to} \quad & \begin{cases} Ax = b \\ x \geq 0. \end{cases} \end{aligned}$$

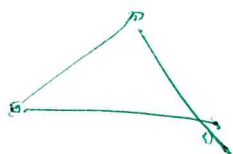
Here we assume that $b \geq 0$ and $\text{rank}(A) = m$. Let the columns of A be given by a_j , i.e. $A = [a_1, a_2, \dots, a_n]$. Let B be an $m \times m$ non-singular matrix whose columns are linearly independent

Simplex



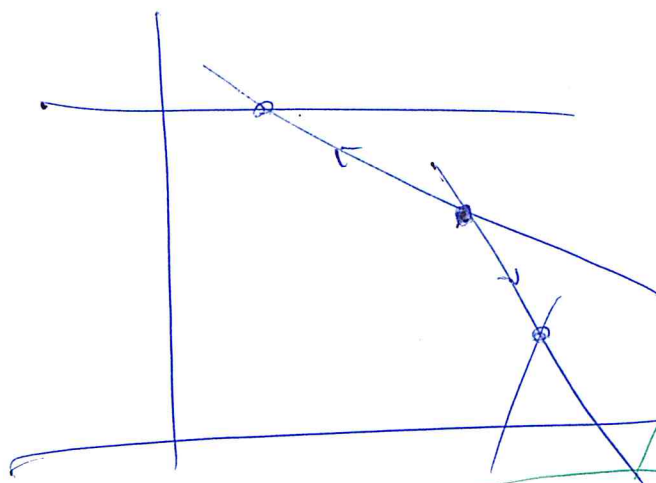
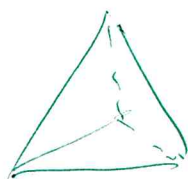
\mathbb{R}^2 3 pts in \mathbb{R}^2 convex hull

Simplex



$n+1$ vertices in \mathbb{R}^n

\mathbb{R}^3 4 pts in \mathbb{R}^3



(i) optimality condition

(ii) feasibility condition

FB \hookrightarrow BFS

Thm 2.2

爬山子法

The Simplex Method for a Two-variable Problem

0.1 Interpretation of the Graphical Method

To introduce the basic ideas of the simplex method, we will use an example with only two decision variables x and y . We can then see how both the graphical method and the simplex method works. Consider

$$\begin{aligned} \text{Max } & f(x, y) = 30x + 20y \\ \text{subject to } & \begin{cases} x + y \leq 50 \\ 40x + 60y \leq 2400 \\ x, y \geq 0 \end{cases} \end{aligned} \quad \text{FCF} \quad (0.1)$$

The graph of the feasible region is given in Figure 0.1.

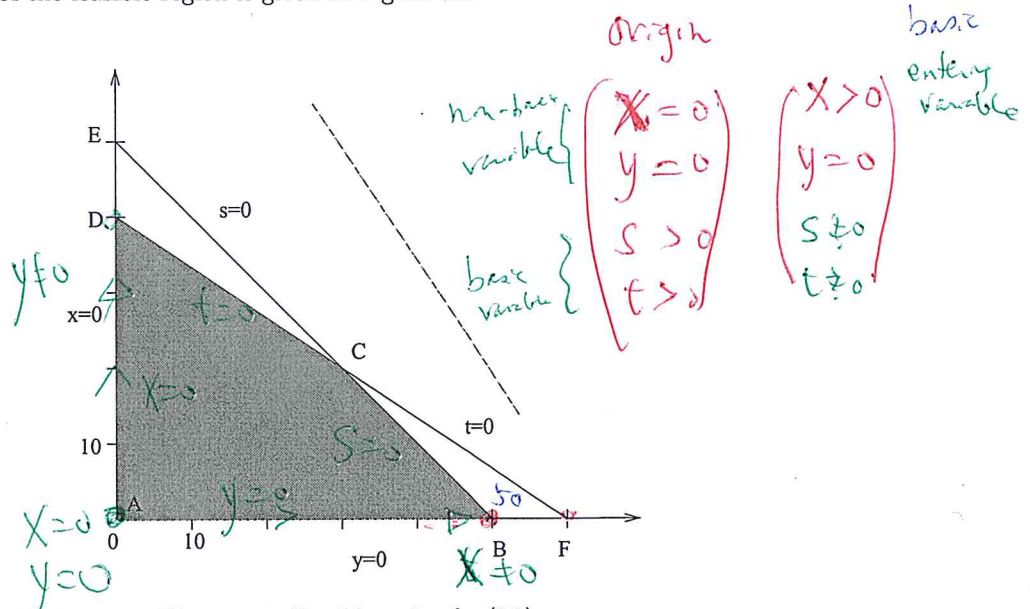
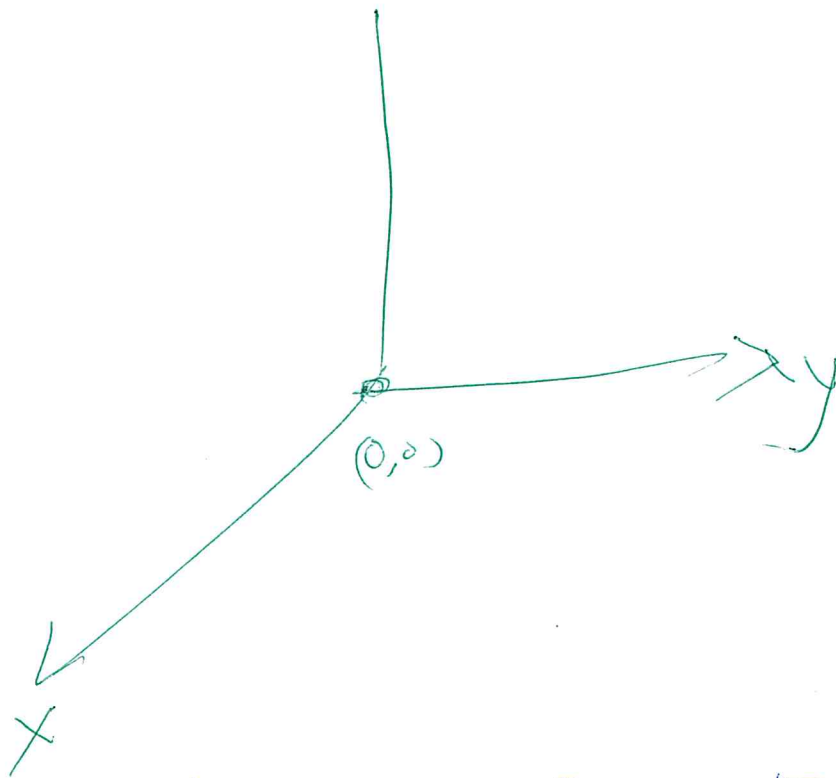


Figure 0.1. Feasible region for (0.1)

We note that the constraints are inequalities. Since inequalities are difficult to be handled by matrices, we first change them into equalities by adding two more variables



$$\begin{array}{r}
 X + 0 + S = 50 \\
 40X \qquad \qquad \qquad + t = 2400
 \end{array}
 \Rightarrow$$

feasibility cut

$$\begin{array}{l}
 X = 50 \\
 X = 60
 \end{array}$$